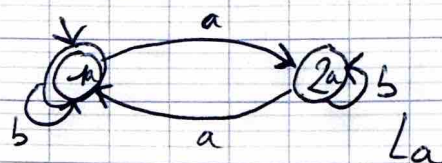


Théorie des langages

Correction du TD n°3

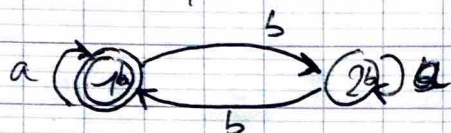
Exercice 1

Automate qui reconnaît un nombre pair de a .
On le note L_a :



$$Q_a = \{1_a, 2_a\}$$

Automate qui reconnaît des mots $|w|_b$ pair

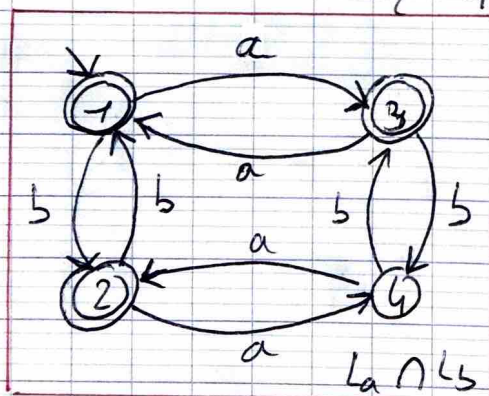


$$Q_b = \{1_b, 2_b\}$$

Construire $L_a \cap L_b$

$$Q_{a \cap b} = Q_a \times Q_b = \{(1_a 1_b), (1_a 2_b), (2_a 1_b), (2_a 2_b)\}$$

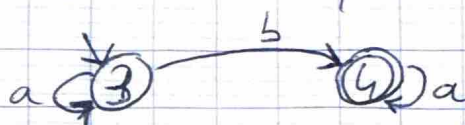
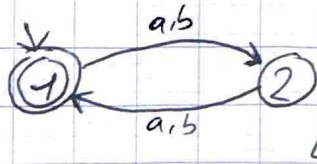
$$= \{1, 2, 3, 4\}$$



Les liens sont effectués par accessibilité dans chaque automates.

Exercice 2

- L_a reconnaît les mots pairs
- L_b reconnaît les mots w tel que $|w|_b \leq 1$

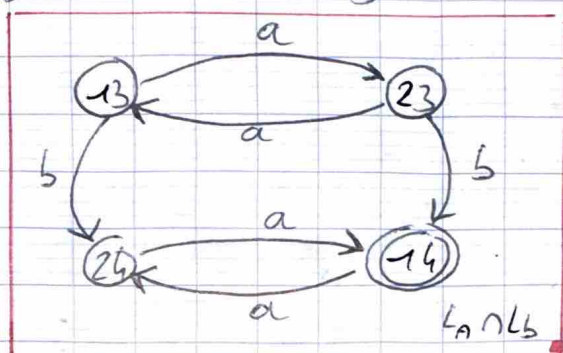


Automate $L_a \cap L_b$.

$$Q_a = \{1, 2\}$$

$$Q_b = \{3, 4\}$$

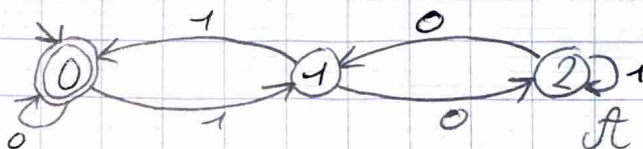
$$Q_a \times Q_b = \{(1,3), (1,4), (2,3), (2,4)\}$$



On privilégie le fait de ne pas créer les arêtes.

Exercice 3 $\Sigma = \{0, 1\}$

On considère l'automate suivant:



Question (a)

$$L(A) = L_0$$

$$\begin{cases} L_0 = 0L_0 + 1L_1 + \epsilon \\ L_1 = 1L_0 + 0L_2 \\ L_2 = 1L_2 + 0L_1 \end{cases}$$

$$L_1 = 1L_0 + 0L_2$$

$$L_2 = 1L_2 + 0L_1$$

$$L_2 = 1L_2 + 0L_1 = 1^* 0L_1$$

$$L_1 = 1L_0 + 01^* 0L_1 = (01^* 0)^* 1L_0$$

$$L_0 = 0L_0 + 1(01^* 0)^* 1L_0 + \epsilon$$

$$= (0 + 1(01^* 0)^* 1)L_0 + \epsilon = (0 + 1(01^* 0)^* 1)^* \epsilon$$

→
Règle d'Anders

Question (b)

les multiples de 3

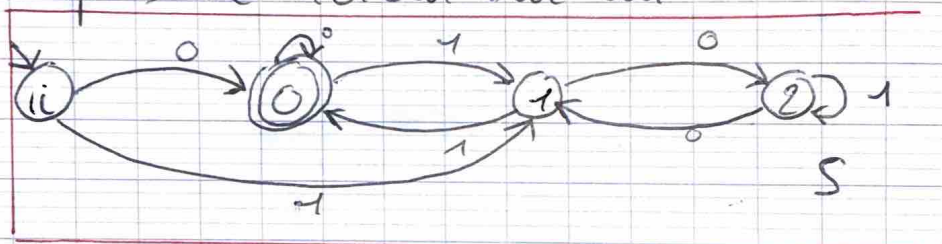
- $0 = (0)_2$
- $3 = (11)_2$
- $6 = (110)_2$
- $9 = (1001)_2$
- $\vdots$

Question (c)

Automate standard

↳ 1 seul état initial

↳ pas de retour sur lui



Question (d)

$$L(S) = L_i$$

$$\begin{cases} L_i = 0L_0 + 1L_1 \\ L_0 = 1L_1 + \epsilon + 0L_0 \\ L_1 = 0L_2 + 1L_0 \\ L_2 = 1L_2 + 0L_1 \end{cases}$$

$$L_2 = 1L_2 + 0L_1 \xrightarrow{A} 1^* 0L_1$$

$$L_1 = 01^* 0L_1 + 1L_0 \xrightarrow{A} (01^* 0)^* 1$$

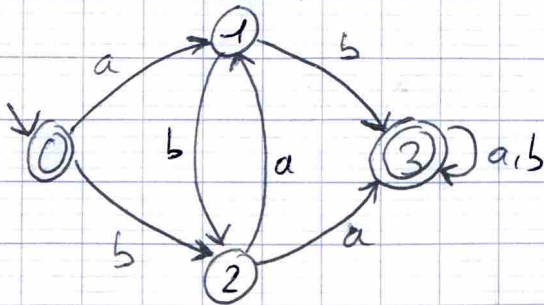
$$L_0 = 1(01^* 0)^* 1 + 0L_0 + \epsilon$$

$$= 0^* 1(01^* 0)^* 1 + \epsilon$$

Exercice 5

L'exercice 4 permet d'évaluer les mêmes compétences.

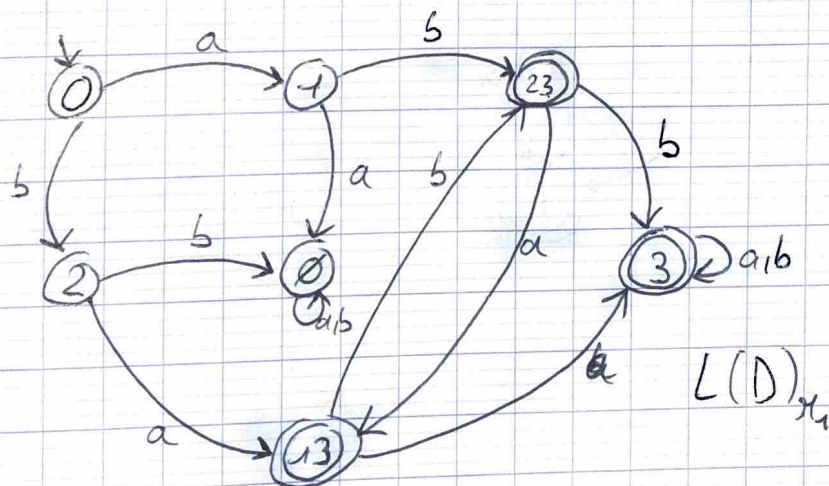
Automate M_1



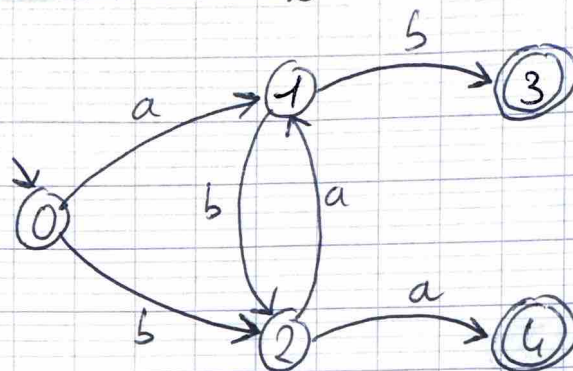
I	∅						
F					∅	∅	∅
	0	1	2	∅	23	13	3
a	1	∅	13	∅	13	3	3
b	2	23	∅	∅	3	23	3

Initial
final

On obtient

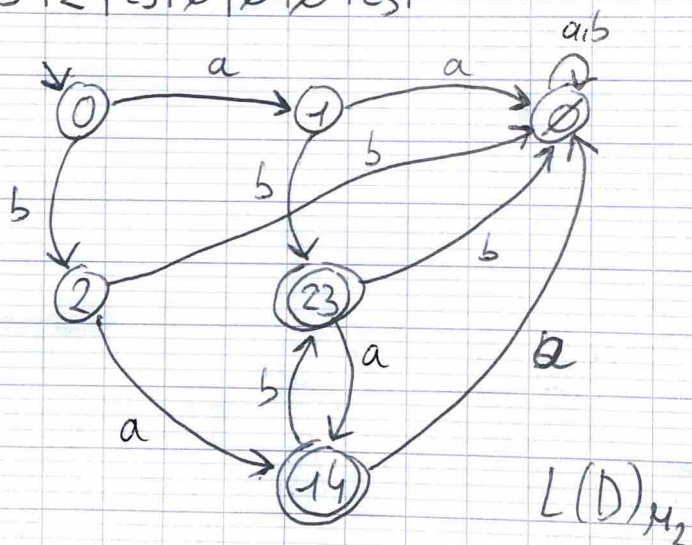


Automate M_2



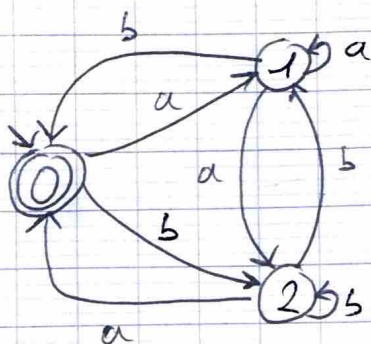
I	∅					
F				∅	∅	∅
	0	1	2	23	∅	14
a	1	∅	14	14	∅	∅
b	2	23	∅	∅	∅	23

Initial
Final



$L(D)_{M_2}$

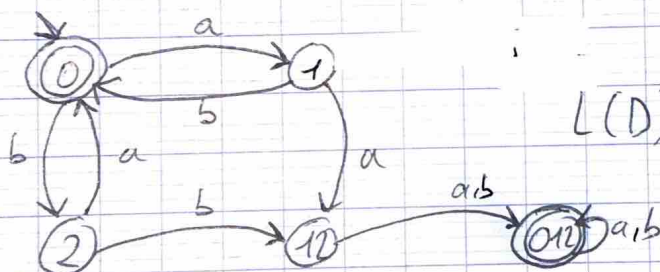
Automate M_3



I	δ				
F	δ				δ
	0	1	2	12	012
a	1	12	0	012	012
b	2	0	12	012	012

Initial

Final



$L(D)_{M_3}$